

CS 70 离散数学 和 概率论

DIS 1A

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归纳 Intro

Natural numbers start at 0, 和 存在 always a next one. 对于 predicates on natural numbers the

Note3

principleofinductionis: $n \in \mathbb{N}, P(n) \equiv P(0) \wedge \forall n, P(n) \Rightarrow P(n+1)$.

即, to 证明 $P(n)$ 对于 natural numbers one proves $P(0)$, the 基础情况,

和 $\forall n, P(n) \Rightarrow$

$P(n+1)$, the 归纳 step. In the 归纳 step, the assumption that

$P(n)$ is 真 is called the

inductionhypothesiswhichistypicallyusedtoarguethat $P(n+1)$ istrue.

An 例子 is the statement $P(n) = \sum_{i=0}^n i = n(n+1)/2$. The 基础情况, $P(0)$,

is the observation that

$\sum_{i=0}^0 i = 0$

Intheinductionstep, theinductionhypothesis, $P(n)$, is $\sum_{i=0}^n i = n(n+1)/2$.

Theinduction

stepproceedsasfollows:

$\sum_{i=0}^{n+1} i = \sum_{i=0}^n i + (n+1) = n(n+1)/2 + (n+1)$

$= (n+1)(n+2)/2$

$= \sum_{i=0}^{n+1} i$

$\sum_{i=0}^{n+1} i = \sum_{i=0}^{n+1} i$

$\sum_{i=0}^{n+1} i = \sum_{i=0}^{n+1} i$

Thefirstequalityfollowsfromthedefinitionofthenotation, \sum , thesecondsubstitutestheinduction hypothesisandthelastisalgebra.

Andwhatisprovenis $P(n+1)$, whichisthat $\sum_{i=0}^{n+1} i = (n+1)(n+2)/2$.

Another 和 等价 view of the natural numbers are that 存在 the

numbers 0 to n 和 那么

thereisn+1. Thestronginductionprincipleisthat

$n \in \mathbb{N}, P(n) \equiv P(0) \wedge \forall n, (\forall k \leq n, P(k)) \Rightarrow P(n+1)$.

Here the 归纳 假设 is that $P(k)$ is 真 对所有 values $k \leq n$. To 证明 that every natural

number $n \geq 2$ can be written as a product of primes, we take

the 基础情况 as $P(2)$ 其中 can be

written as 2, 其中 is a product of a prime. 和 对于 any n, 如果 it

is prime, it can be written as itself,

otherwise $n=ab$ 和 by the inductive hypotheses $P(a)$ 和 $P(b)$ is

that each can be written as a

product of primes. 因此, we can write n as the product of the

primes in both a 和 b. 笔记 here

that the 基础情况 starts at 2, 其中 illustrates that one chose the

基础情况 as is relevant to the

statementbeingproven.

Strengtheningtheinductionhypothesisisatechniquethatprovesastrongertheorem.

Forexample,

the notes 考虑 the 定理 "The sum of the first n odd numbers is a perfect square." In fact,

the notes inductively 证明 the stronger 定理 "The sum of the first n odd numbers is n^2 ." Here,

the stronger inductive 假设 allows the 归纳 step to proceed

easily. 笔记 that in strong

归纳, we 假设 more cases are 真 in the inductive 假设, whereas strengthening the

inductive hypothesis proves a stronger claim entirely.
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1 Natural 归纳 on Inequality

Prove that if $n \in \mathbb{N}$ and $x > 0$, then $(1+x)^n \geq 1+nx$.

Note 3

解答:

- Base Case: When $n=0$, the claim holds since $(1+x)^0 \geq 1+0x$.

- Inductive Hypothesis: Assume that $(1+x)^k \geq 1+kx$

for some value of $n=k$ where $k \in \mathbb{N}$.

- Inductive Step: For $n=k+1$, we can show the following:

$$(1+x)^{k+1} = (1+x)^k (1+x) \geq (1+kx)(1+x)$$

$$\geq 1+kx+x+kx^2$$

$$\geq 1+(k+1)x+kx^2 \geq 1+(k+1)x$$

By induction, we have shown that $n \in \mathbb{N}$, $(1+x)^n \geq 1+nx$.

2 Make It Stronger

假设 that the 序列 a_1, a_2, \dots is defined by $a_1 = 1$ 和 $a_n = 3a_{n-1}$ 对于 $n \geq 2$.

We want to

Note 3 $1 \leq a_n \leq 3^{n-1}$

prove that

$$a_n \leq 3^{2n}$$

for every positive integer n .

(a) 假设 that we want to 证明 this statement using 归纳. Can we 令 our inductive

假设 be simply $a_n \leq$

$$3^{2n}?$$

Attempt an 归纳 证明 with this 假设 to 证明

why this does not work.

(b)

Instead, 令' s try to instead prove a stronger statement instead, of the form

$$a_n \leq 3^{2n} \text{ ??}.$$

What

should replace the 问题 marks 所以 that the 归纳步骤 works out? Try some examples

and see what works.

(c)

Using your stronger inductive hypothesis from part (b), carry out the induction proof.

(d) Why does the hypothesis in part (b) imply the overall claim?

解答:

(a) 令' s try to prove that for every $n \geq 1$, we have

$$a_n \leq 3^{2n}$$

by induction.

n

Base Case: For $n=1$ we have

$$a_1 = 1 \leq 3^{2 \cdot 1}$$

$$= 9.$$

n

Inductive Step: For some $n \geq 1$, we assume

$$a_n \leq 3^{2n}$$

. Now, consider $n+1$. We can write:

n

$$a_{n+1} = 3a_n \leq 3(3^{2n})^2 = 3 \times 3^{2 \times 2n} = 3 \times 3^{4n} = 3^{4n+1} = 3^{2(n+1)}.$$

$n+1$

然而, what we wanted was to get an inequality of the form: $a_n \leq$

$$3^{2n+1}$$

. 存在 an

n+1

extra+1 in the exponent of what we derived.

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(b) The 假设 should be $a \leq 3 \cdot 2^{n-1}$.

$3 \cdot 2^{n-1}$.

We can see in the next part that this modification

n

allows the algebra in the inductive step to work out nicely.

(c) This time the induction works.

Base Case: For $n=1$ we have $a=1 \leq 3 \cdot 2^{1-1} = 3$.

1

Inductive Step: For some $n \geq 1$ we assume $a \leq 3 \cdot 2^{n-1}$.

$\leq 3 \cdot 2^{n-1}$.

Now, consider $n+1$. We can write:

n

$a = 3a \leq 3 \times (3 \cdot 2^{n-1})^2 = 3 \times 3^2 \times (2^{n-1})^2 = 3 \times 3^2 \times 2^{2n-2} = 3 \cdot 2^{2n-1}$.

n+1 n

This is exactly the induction hypothesis for $n+1$.

(d)

For every $n \geq 1$, we have $2^{n-1} \leq 2^n$ and therefore $3 \cdot 2^{n-1} \leq 3 \cdot 2^n$.

. This means that our modified

假设 其中 we proved in part (b) does indeed imply what we wanted

to 证明 in part

(a).

3 Binary Numbers

证明 that every 正的 整数 n can be written in binary. 换句话说, 证明

that 对于 any

Note 3

positive integer n, we can write

$n = c_k \cdot 2^k + c_{k-1} \cdot 2^{k-1} + \dots + c_1 \cdot 2^1 + c_0 \cdot 2^0$,

$c_i \in \{0, 1\}$

for some $k \in \mathbb{N}$ and $c_i \in \{0, 1\}$ for all $i \leq k$.

i

As an 例子, the number 13 can be written in binary as 1101

because $13 = 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0$.

$2^1 + 1 \cdot 2^0$.

解答:

Prove by strong induction on n.

The key insight here is that 如果 n is divisible by 2, 那么 it

is easy to get a bit string representation

of (n+1) from that of n. 然而, 如果 n is 非 divisible by 2, 那么

(n+1) will be, 和 its binary

representation will be more easily derived from that of (n+1)/2.

More formally:

• Base Case: $n=1$ can be written as 1×2^0 .

• Inductive Step: 假设 that the statement is 真 对所有 $1 \leq m \leq n$, 其中 n is arbitrary.

Now, we need to 考虑 n+1. 如果 n+1 is divisible by 2, 那么 we can apply our inductive

hypothesis to $(n+1)/2$ and use its representation to express $n+1$ in the desired form.

$(n+1)/2 = c_k \cdot 2^k + c_{k-1} \cdot 2^{k-1} + \dots + c_1 \cdot 2^1 + c_0 \cdot 2^0$

$c_i \in \{0, 1\}$

$n+1 = 2 \cdot (n+1)/2 = c_k \cdot 2^{k+1} + c_{k-1} \cdot 2^k + \dots + c_1 \cdot 2^2 + c_0 \cdot 2^1 + 0 \cdot 2^0$.

$c_i \in \{0, 1\}$

Otherwise, n must be divisible by 2 和 因此 have $c=0$. We can obtain the representation

0

of $n+1$ from n as follows:

$n = c \cdot 2^{k+c} \cdot 2^{k-1} + \dots + c \cdot 2^{1+0} \cdot 2^0$
 $k \quad k-1 \quad 1$
 $n+1 = c \cdot 2^{k+c} \cdot 2^{k-1} + \dots + c \cdot 2^{1+1} \cdot 2^0$
 $k \quad k-1 \quad 1$
 因此, the statement is true.
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Here is another alternate 解答 emulating the 算法 of converting a decimal number to a binary number.

- Base Case: $n=1$ can be written as 1×2^0 .
- Inductive Step: Assume that the statement is true for all $1 \leq m \leq n$, for arbitrary n . We show that the statement holds for $n+1$. Let 2^m be the largest power of 2 such that $n+1 \geq 2^m$. 因此, $n+1 < 2^{m+1}$. We examine the number $(n+1) - 2^m$. 因为 $(n+1) - 2^m < n+1$, the inductive hypothesis holds, 所以我们有 a binary representation 对于 $(n+1) - 2^m$. (如果 $(n+1) - 2^m = 0$, then we still have a binary representation, namely $0 \cdot 2^0$.) Also, 因为 $n+1 < 2^{m+1}$, $(n+1) - 2^m < 2^m$, 所以 the largest power of 2 in the representation of $(n+1) - 2^m$ is 2^{m-1} . 因此, by the inductive hypothesis, $(n+1) - 2^m = c \cdot 2^{m-1+c} \cdot 2^{m-2} + \dots + c \cdot 2^{1+c} \cdot 2^0$, and adding 2^m to both sides gives $n+1 = 2^m + c \cdot 2^{m-1+c} \cdot 2^{m-2} + \dots + c \cdot 2^{1+c} \cdot 2^0$, which is a binary representation for $n+1$. 因此, the induction is complete.

Another intuition is that if x has a binary representation, $2x$ and $2x+1$ do as well: shift the bits and possibly place 1 in the last bit. The above 归纳 could 那么 have proceeded from n 和 used the binary representation of $n/2$, shifting 和 possibly setting the first bit depending on whether n is odd or even.

笔记: In proofs using simple 归纳, we only use $P(n)$ in order to 证明 $P(n+1)$. Simple 归纳 gets stuck here because in order to 证明 $P(n+1)$ in the 归纳步骤, we need to assume more than just $P(n)$.

This is because it is not immediately clear how to get a representation 对于 $P(n+1)$ using just $P(n)$, particularly in the case that $n+1$ is divisible by 2. As a result, we

assume the statement to be true for all of $1, 2, \dots, n$ in order to prove it for $P(n+1)$.

4 Fibonacci 对于 Home
 Recall, the Fibonacci numbers, defined recursively as
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$F_1 = 1, F_2 = 1$, and $F_n = F_{n-1} + F_{n-2}$.

Prove that every third Fibonacci number is even. For example, $F_3 = 2$ is even and $F_6 = 8$ is even.

3 6
 解答:

We want to prove that for all natural numbers $k \geq 1$, F_{3k} is even.

Basecase: For $k=1$, we can see that $F_1 = 1$ is even.

3

Induction hypothesis:

Suppose that for an arbitrary fixed value of k , F_k is even.

3k

Inductive step: We can write

$F_{k+1} = F_k + F_{k-1} = 2F_k + F_{k-1}$.

$3k+3 \quad 3k+2 \quad 3k+1 \quad 3k+1 \quad 3k$

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By the induction hypothesis, we know that $F_k = 2q$ for some q .

3k

This means that 我们有 that $F_{k+1} = 2(2q + F_{k-1})$, 其中 蕴含 that it is even.

因此, by the

$3k+3 \quad 3k+1$

principles of induction we have shown that all F_k are even.

3k

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